## An introduction to persistent homology

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## Part I: Scalar Functions

Persistence Diagram of a scalar function

- Track the evolution of the topology of sub-level sets as the threshold increases.
- Pair thresholds that create components with those that destroy them.



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## Persistence Diagram of a scalar function

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Algorithm 3: Calculating discrete o-dimensional persistent homology
Require: A discrete sample \(\left\{\left(x_{1}, y_{1}\right),\left(x_{2}, y_{2}\right), \ldots\right\}\) of a function \(f: \mathbb{D} \subseteq \mathbb{R} \rightarrow \mathbb{R}\)
    function PersistentHomology \((f)\)
        \(\mathrm{U} \leftarrow \varnothing \quad \triangleright\) Initialize an empty union-find structure
        Sort the value tuples in ascending order, such that \(y_{1} \geq y_{2} \geq \ldots\)
        for Tuple \(\left(x_{i}, y_{i}\right)\) of \(f\) do
            if \(y_{i-1}>y_{i}\) and \(y_{i+1}>y_{i}\) then \(\quad \triangleright y_{i}\) is a local minimum
                    U.add(i)
            else if \(y_{i-1}<y_{i}\) and \(y_{i+1}<y_{i}\) then \(\quad \triangleright y_{i}\) is a local maximum
            \(c \leftarrow \operatorname{U} \cdot \operatorname{get}(i-1)\)
            \(d \leftarrow \operatorname{U.get}(i+1) \quad \triangleright\) Get second connected component
            U.merge \((c, d) \quad \triangleright\) Merge the two connected components meeting at \(y_{i}\)
            else
            \(c \leftarrow \operatorname{U} . \operatorname{get}(i-1)\)
            \(\mathrm{U}[c] \leftarrow \mathrm{U}[c] \cup i \quad \triangleright\) Add \(y_{i}\) to the current connected component
            end if
        end for
        return U
    end function
```


## The pairing algorithm

- Input : a discrete sample $P=\left\{p_{1}=\left(x_{1}, y_{1}\right), \ldots, p_{n}=\left(x_{n}, y_{n}\right)\right\}$ representing a scalar function f .
- A collection of paired points.

1. Initiate an empty UnionFind $U$.
2. Sort $P$ with respect the $y$ values.
3. For every $p_{i}=\left(x_{i}, y_{i}\right)$ in $P$ :
4. If $y_{i-1}>y_{i}$ and $y_{i+1}>y_{i}$ then:
5. U.add(i)
6. Set the birth of i to $y_{i}$
7. Else if $y_{i-1}<y_{i}$ and $y_{i+1}<y_{i}$ then:
8. $\mathrm{c}=\mathrm{U} . \operatorname{get}(\mathrm{i}-1)$
9. $d=U . \operatorname{get}(i+1)$
10. U.merge(c,d)
11. Pair $i$ with c or d (choose the one that was born later)
12. Else:
13. $c=U . \operatorname{get}(i-1)$
14. $\mathrm{U}(\mathrm{c}):=\mathrm{U}(\mathrm{c})$ union i

## Part II : Point Clouds Introduction to VR and Cech complexes

## Nerve of a topological space



Given a set of points P sampled from a space X , how can we recover the topological features of the original space $X$ from the point cloud $P$ ?

## Nerve of a topological space



We want a discretized structure that capture the shape of the space and we want a reasonable way that is subtle enough to measure this shape.

Nerve of a topological space


## Čech complex

Given a point cloud X in some metric space and a number $\varepsilon>0$, the Čech complex $C_{\varepsilon}$ is the simplicial complex whose simplices are constructed as follows :

For each subset Y of X , form a $(\varepsilon / 2)$-ball around each point in Y , and include Y as a simplex ,of dimension $|\mathrm{Y}|$, if there is a common point contained in all of the balls in Y .


## The Cech complex approximates the topological space

Theorem: The homotopy type of $S_{\epsilon}$ and $C_{\varepsilon}$ are the same.


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## Čech complex size

For each subset Y of X , form a $(\varepsilon / 2)$-ball around each point in Y , and include Y as a simplex ,of dimension $|\mathrm{Y}|$, if there is a common point contained in all of the balls in Y .

What is the computational problem in constructing a Čech complex?

If we have a point cloud set $X$ of size 40 then we have to check all subsets of $X$ of size 40 . This is $2^{\wedge}\{40\}$. Very slow!

## Vietoris-Rips complex

Let $X$ is a subset of a metric space $d$ and let $\epsilon>0$. The Vietoris-Rips complex is constructed as follows :
(1) For each point in $X$, make it as a 0 -simplex.
(2) For each pair $x_{1}, x_{2} \in X$, make a 1 -simplex $\left(\left[x_{1}, x_{2}\right]\right)$ if $d\left(x_{1}, x_{2}\right) \leq \epsilon$.
(3) For $x_{1}, x_{2}, \cdots, x_{n} \in X$, make an $(n-1)$-simplex with vertices $x_{1}, x_{2}, \cdots, x_{n}$. Then, $d(x i, x j) \leq \epsilon$ for all $0 \leq i, j \leq n$; that is, if all the points are within a distance of $\epsilon$ from each other.

This complex is denoted by $\operatorname{VR}(X, \epsilon)$


## Čech complex and VR complex

Comparison between the two complexes :


Note that the VR complex does not necessarily have the same homotopy type of the space of the union of ball.

## Čech complex and VR complex

What is the relation between the Čech complex and VR complex ?

Theorem: For all $\varepsilon>0$, the following inclusions hold

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C_{\varepsilon} \subset V R_{\varepsilon} \subset C_{2 \varepsilon}
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So the VR complex forms a good approximation of the Čech complex.

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- The number of connected components,
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- The Number o voids in a space


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This space here has 1 connected component and 3 cycles

## What size do we consider?



We will come back
to this question
later.

## Remarks on computing the Vietoris-Rips complex

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Let $\varepsilon_{1} \leq \varepsilon_{2}$ then $V R\left(X, \varepsilon_{1}\right) \leq V R\left(X, \varepsilon_{2}\right)$

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This allows us in practice to compute the VR complex for some maximum scale a $\in R$ and then extract the complex at any lower scale b less than a.

Given a $\operatorname{VR}(X, a)$, suppose that we want to compute $\operatorname{VR}(X, b)$ for b less than a. How do we compute determine the simplices from $\operatorname{VR}(\mathrm{X}, \mathrm{a})$ that belongs to $\operatorname{VR}(\mathrm{X}, \mathrm{b})$ ?

## Remarks on computing the Vietoris-Rips complex

Given a complex $V R(X, \varepsilon)$ define the weight function $w: V R(X, \varepsilon) \rightarrow \mathbb{R}$

$$
\omega(\sigma)= \begin{cases}0, & \operatorname{dim}(\sigma) \leq 0 \\ \mathrm{~d}(u, v), & \sigma=\{u, v\} \\ \max _{\tau \subset \sigma} \omega(\tau), & \text { otherwise }\end{cases}
$$

That is, the weight $\omega(\sigma)$ is equal to the maximum of the weights (lengths) of all its edges.

After defining the weight function on $V R\left(X, \varepsilon_{2}\right)$ we sort the simplices according to their weights, extracting the VR complex for any $\varepsilon_{1} \leq \varepsilon_{2}$ as a prefix of this ordering.

This gives a filtration of $\operatorname{VR}\left(X, \varepsilon_{2}\right)$

## The relation between neighborhood graph and the Vietoris-Rips complex



The data (left) has the $\boldsymbol{\epsilon}$-neighborhood graph (middle).
This is precisely the VR complex (right) at that same resolution.

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Question: given the $\epsilon$-neighborhood (middle), how can we recover the VR complex from it (right) ?

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This is precisely the VR complex (right) at that same resolution.
Question: given the $\epsilon$-neighborhood (middle), how can we recover the VR complex from it (right) ?
Answer: Higher dimensional simplices recovered from the cliques of the $\boldsymbol{\epsilon}$-neighborhood graph.

