# An Introduction to Multidimensional Scaling and ISOMAP 

## MUSTAFA HAJIJ

## MDS

Let $D=\left[d_{i j}\right]$ be an $N \times N$ dissimalrity matrix.

In MDS, we want to find $n$ vectors $x_{1}, \ldots, x_{N}$ in $R^{d}$ such that $\left\|x_{i}-x_{j}\right\| \approx d_{i j}$

## MDS

Let $D=\left[d_{i j}\right]$ be an $N \times N$ dissimalrity matrix.

In MDS, we want to find $n$ vectors $x_{1}, \ldots, x_{N}$ in $R^{d}$ such that $\left\|x_{i}-x_{j}\right\| \approx d_{i j}$

- Usually if we choose d to be large enough, we can construct the vectors $x_{1}, \ldots, x_{N}$ with exact solutions : $\left\|x_{i}-x_{j}\right\|=d_{i j}$.


## MDS

Let $D=\left[d_{i j}\right]$ be an $N \times N$ dissimalrity matrix.

In MDS, we want to find $n$ vectors $x_{1}, \ldots, x_{N}$ in $R^{d}$ such that $\left\|x_{i}-x_{j}\right\| \approx d_{i j}$

- Usually if we choose d to be large enough, we can construct the vectors $x_{1}, \ldots, x_{N}$ with exact solutions : $\left\|x_{i}-x_{j}\right\|=d_{i j}$.
- In this case the distance $d$ above is the usual Euclidean distance.


## MDS

Let $D=\left[d_{i j}\right]$ be an $N \times N$ dissimalrity matrix.

In MDS, we want to find $n$ vectors $x_{1}, \ldots, x_{N}$ in $R^{d}$ such that $\left\|x_{i}-x_{j}\right\| \approx d_{i j}$

- Usually if we choose d to be large enough, we can construct the vectors $x_{1}, \ldots, x_{N}$ with exact solutions : $\left\|x_{i}-x_{j}\right\|=d_{i j}$.
- In this case the distance $d$ above is the usual Euclidean distance.
- There are cases where the matrix D is valid distance matrix, but still there exists no set of vectors $x_{1}, \ldots, x_{N}$ in any $R^{d}$ with perfect $\left\|x_{i}-x_{j}\right\|=d_{i j}$. Such a distance is called nonEuclidean distance.


## Classical MDS Algorithm

1-Construct the matrix of squares of the distances $P^{(2)}=\left[d_{i j}^{2}\right]$.

## Classical MDS Algorithm

1-Construct the matrix of squares of the distances $P^{(2)}=\left[d_{i j}^{2}\right]$.
2. Apply the double centering: $\mathrm{B}=-\frac{1}{2} J P^{(2)} J$ where $J=I-\frac{1}{n} 11^{\prime}$, where N is the number of elements.

## Classical MDS Algorithm

1-Construct the matrix of squares of the distances $P^{(2)}=\left[d_{i j}^{2}\right]$.
2. Apply the double centering: $\mathrm{B}=-\frac{1}{2} J P^{(2)} J$ where $J=I-\frac{1}{n} 11^{\prime}$, where N is the number of elements.
3. Extract the largest d positive eigenvalues $\lambda_{1} \ldots \lambda_{d}$ of $B$ and the corresponding m eigenvectors $e_{1} \ldots e_{d}$.

## Classical MDS Algorithm

1-Construct the matrix of squares of the distances $P^{(2)}=\left[d_{i j}^{2}\right]$.
2. Apply the double centering: $\mathrm{B}=-\frac{1}{2} J P^{(2)} J$ where $J=I-\frac{1}{n} 11^{\prime}$, where N is the number of elements.
3. Extract the largest d positive eigenvalues $\lambda_{1} \ldots \lambda_{d}$ of $B$ and the corresponding $m$ eigenvectors $e_{1} \ldots e_{d}$.
4. A d-dimensional MDS coordinates of n objects is derived from the coordinate matrix $X=E_{d} \Lambda_{\mathrm{d}}^{\frac{1}{2}}$, where $E_{d}$ is the matrix of d eigenvectors and $\Lambda_{d}$ is the diagonal matrix of $d$ eigenvalues of $B$, respectively

## Example

Start with a distance matrix D

$\mathrm{D}=$| 0 | 93 | 82 | 133 |
| ---: | ---: | ---: | ---: |
| 93 | 0 | 52 | 60 |
| 82 | 52 | 0 | 111 |
| 133 | 60 | 111 | 0 |

## Example

Take the square the elements of $D$

$\mathbf{D}=$| 0 | 93 | 82 | 133 |
| ---: | ---: | ---: | ---: |
| 93 | 0 | 52 | 60 |
| 82 | 52 | 0 | 111 |
| 133 | 60 | 111 | 0 |\(\quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrr}0 \& 8649 \& 6724 \& 17689 <br>

8649 \& 0 \& 2704 \& 3600 <br>
6724 \& 2704 \& 0 \& 12321 <br>
17689 \& 3600 \& 12321 \& 0\end{array}\right]\)

$$
\begin{gathered}
\mathbf{D}=\begin{array}{rrrr}
0 & 93 & 82 & 133 \\
93 & 0 & 52 & 60 \\
82 & 52 & 0 & 111 \\
133 & 60 & 111 & 0
\end{array} \quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrr}
0 & 8649 & 6724 & 17689 \\
8649 & 0 & 2704 & 3600 \\
6724 & 2704 & 0 & 12321 \\
17689 & 3600 & 12321 & 0
\end{array}\right] \\
\mathbf{J}=\left[\begin{array}{lrrr}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right]-0.25 \times\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1
\end{array}\right]=\left[\begin{array}{rrrr}
0.75 & -0.25 & -0.25 & -0.25 \\
-0.25 & 0.75 & -0.25 & -0.25 \\
-0.25 & -0.25 & 0.75 & -0.25 \\
-0.25 & -0.25 & -0.25 & 0.75
\end{array}\right]
\end{gathered}
$$

$\mathrm{D}=\begin{array}{rrrr}0 & 93 & 82 & 133 \\ 93 & 0 & 52 & 60 \\ 82 & 52 & 0 & 111 \\ 133 & 60 & 111 & 0\end{array} \quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrr}0 & 8649 & 6724 & 17689 \\ 8649 & 0 & 2704 & 3600 \\ 6724 & 2704 & 0 & 12321 \\ 17689 & 3600 & 12321 & 0\end{array}\right]$
$\mathbf{J}=\left[\begin{array}{llll}1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1\end{array}\right]-0.25 \times\left[\begin{array}{llll}1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1\end{array}\right]=\left[\begin{array}{rrrr}0.75 & -0.25 & -0.25 & -0.25 \\ -0.25 & 0.75 & -0.25 & -0.25 \\ -0.25 & -0.25 & 0.75 & -0.25 \\ -0.25 & -0.25 & -0.25 & 0.75\end{array}\right]$
$\mathbf{B}=-\frac{1}{2} \mathbf{J P}^{(\mathbf{2})} \mathbf{J}=\left[\begin{array}{rrrr}5035.0625 & -1553.0625 & 258.9375 & -3740.938 \\ -1553.0625 & 507.8125 & 5.3125 & 1039.938 \\ 258.9375 & 5.3125 & 2206.8125 & -2471.062 \\ -3740.9375 & 1039.9375 & -2471.0625 & 5172.062\end{array}\right]$

$$
\begin{gathered}
\left.\mathbf{D}=\begin{array}{rrrr}
0 & 93 & 82 & 133 \\
93 & 0 & 52 & 60 \\
82 & 52 & 0 & 111 \\
133 & 60 & 111 & 0
\end{array}\right] \quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrrr}
0 & 8649 & 6724 & 17689 \\
8649 & 0 & 2704 & 3600 \\
6724 & 2704 & 0 & 12321 \\
17689 & 3600 & 12321 & 0
\end{array}\right] \\
\mathbf{J}=\left[\begin{array}{lrrr}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right]-0.25 \times\left[\begin{array}{lrrr}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1
\end{array}\right]=\left[\begin{array}{rrrr}
0.75 & -0.25 & -0.25 & -0.25 \\
-0.25 & 0.75 & -0.25 & -0.25 \\
-0.25 & -0.25 & 0.75 & -0.25 \\
-0.25 & -0.25 & -0.25 & 0.75
\end{array}\right] \\
\mathbf{B}=-\frac{1}{2} \mathbf{J P}^{(\mathbf{2})} \mathbf{J}=\left[\begin{array}{rrrr}
5035.0625 & -1553.0625 & 258.9375 & -3740.938 \\
-1553.0625 & 507.8125 \\
258.9375 & 5.3125 & 2206.8125 & 1039.938 \\
-3740.9375 & 1039.9375 & -2471.0625 & 5172.062
\end{array}\right]
\end{gathered}
$$

$\lambda_{1}=9724.168, \lambda_{2}=3160.986, \quad \mathbf{e}_{1}=\left(\begin{array}{r}-0.637 \\ 0.187 \\ -0.253 \\ 0.704\end{array}\right) \quad, \quad \mathbf{e}_{2}=\left(\begin{array}{r}-0.586 \\ 0.214 \\ 0.706 \\ -0.334\end{array}\right)$

$$
\begin{gathered}
\left.\mathbf{D}=\begin{array}{rrrr}
0 & 93 & 82 & 133 \\
93 & 0 & 52 & 60 \\
82 & 52 & 0 & 111 \\
133 & 60 & 111 & 0
\end{array}\right] \quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrrr}
0 & 8649 & 6724 & 17689 \\
8649 & 0 & 2704 & 3600 \\
6724 & 2704 & 0 & 12321 \\
17689 & 3600 & 12321 & 0
\end{array}\right] \\
\mathbf{J}=\left[\begin{array}{lrrr}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right]-0.25 \times\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1
\end{array}\right]=\left[\begin{array}{rrrr}
0.75 & -0.25 & -0.25 & -0.25 \\
-0.25 & 0.75 & -0.25 & -0.25 \\
-0.25 & -0.25 & 0.75 & -0.25 \\
-0.25 & -0.25 & -0.25 & 0.75
\end{array}\right] \\
\mathbf{B}=-\frac{1}{2} \mathbf{J P}^{(\mathbf{2})} \mathbf{J}=\left[\begin{array}{rrrr}
5035.0625 & -1553.0625 & 258.9375 & -3740.938 \\
-1553.0625 & 507.8125 & 5.3125 & 1039.938 \\
258.9375 & 5.3125 & 2206.8125 & -2471.062 \\
-3740.9375 & 1039.9375 & -2471.0625 & 5172.062
\end{array}\right]
\end{gathered}
$$

$\lambda_{1}=9724.168, \lambda_{2}=3160.986, \quad \mathbf{e}_{\mathbf{1}}=\left(\begin{array}{r}-0.637 \\ 0.187 \\ -0.253 \\ 0.704\end{array}\right), \mathbf{e}_{\mathbf{2}}=\left(\begin{array}{r}-0.586 \\ 0.214 \\ 0.706 \\ -0.334\end{array}\right) \quad \mathbf{X}=\left[\begin{array}{rr}-0.637 & -0.586 \\ 0.187 & 0.214 \\ -0.253 & 0.706 \\ 0.704 & -0.334\end{array}\right]\left[\begin{array}{r}-62.831 \\ 18.403 \\ -32.97448 \\ -24.960 \\ 39.7109697 \\ 39.388 \\ -18.76340\end{array}\right]$

$$
\begin{gathered}
\left.\mathbf{D}=\begin{array}{rrrr}
0 & 93 & 82 & 133 \\
93 & 0 & 52 & 60 \\
82 & 52 & 0 & 111 \\
133 & 60 & 111 & 0
\end{array}\right] \quad \mathbf{P}^{(\mathbf{2})}=\left[\begin{array}{rrrrr}
0 & 8649 & 6724 & 17689 \\
8649 & 0 & 2704 & 3600 \\
6724 & 2704 & 0 & 12321 \\
17689 & 3600 & 12321 & 0
\end{array}\right] \\
\mathbf{J}=\left[\begin{array}{lrrr}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right]-0.25 \times\left[\begin{array}{lrrr}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1
\end{array}\right]=\left[\begin{array}{rrrr}
0.75 & -0.25 & -0.25 & -0.25 \\
-0.25 & 0.75 & -0.25 & -0.25 \\
-0.25 & -0.25 & 0.75 & -0.25 \\
-0.25 & -0.25 & -0.25 & 0.75
\end{array}\right] \\
\mathbf{B}=-\frac{1}{2} \mathbf{J P}^{(\mathbf{2})} \mathbf{J}=\left[\begin{array}{rrrr}
5035.0625 & -1553.0625 & 258.9375 & -3740.938 \\
-1553.0625 & 507.8125 & 5.3125 & 1039.938 \\
258.9375 & 5.3125 & 2206.8125 & -2471.062 \\
-3740.9375 & 1039.9375 & -2471.0625 & 5172.062
\end{array}\right]
\end{gathered}
$$

Coordinates of first point
$\lambda_{1}=9724.168, \lambda_{2}=3160.986, \quad \mathbf{e}_{\mathbf{1}}=\left(\begin{array}{r}-0.637 \\ 0.187 \\ -0.253 \\ 0.704\end{array}\right), \mathbf{e}_{\mathbf{2}}=\left(\begin{array}{r}-0.586 \\ 0.214 \\ 0.706 \\ -0.334\end{array}\right) \quad \mathbf{X}=\left[\begin{array}{rr}-0.637 & -0.586 \\ 0.187 & 0.214 \\ -0.253 & 0.706 \\ 0.704 & -0.334\end{array}\right]\left[\begin{array}{rr}\sqrt{9724.168} & 0 \\ 0 & \sqrt{3160.986}\end{array}\right]=\left[\begin{array}{rr}-62.831 & -32.97448 \\ -24.960 \\ 69.388 & 12.02697 \\ 39.71091 \\ -18.76340\end{array}\right]$

## Stress Majorization

In the classical MDS algorithm the cost function that we are trying to optimize is called the stress function and it is given by :

$$
\operatorname{Stress}_{D}\left(x_{1}, x_{2}, \ldots, x_{N}\right)=\left(\sum_{i \neq j=1, \ldots, N}\left(d_{i j}-\left\|x_{i}-x_{j}\right\|\right)^{2}\right)^{1 / 2}
$$

In this function we try to find $x_{1}, \ldots, x_{N}$ in a certain dimension $d$ such that $\operatorname{Stress}\left(x_{1}, \ldots, x_{N}\right)$ is as small as possible

## Stress Majorization

The stress function has a more general form as :

$$
\sigma(X)=\sum_{i<j \leq n} w_{i j}\left(d_{i j}(X)-\delta_{i j}\right)^{2}
$$

Here wij weight between a pair of points ( $\mathrm{i}, \mathrm{j}$ ) that represents the confidence in in the similarity between points (i,j).
$\delta_{i j}$ the given distance between the points $\mathrm{i}, \mathrm{j}$

## Stress Majorization

"Pressing" the data into 2 dimensions enables us to visualize the data. However, that comes with a price : high stress function value (which correlates wit distorted representation )


## MDS in Sklearn

## MDS is implemented in Sklearn



## MDS in Sklearn

MDS is implemented in Sklearn


Example

## MDS in Sklearn

MDS is implemented in Sklearn

```
A selection from the 64-dimensional digits dataset
    012%3450112345012234505
55041351002220123333
44150522001321471314
14405315442225544001
23450123450123450555
04135100222042333344
15052200132131314314
O5745441225544001234
50123450423450555041
35100222012333344150
52200132143131431405
34544222554403012345
01234501234505550413
51001220123933441505
22001321431314314053
15442225544001234601
2345012 24450555041354
0022201423333441550522
44222554400123450123
```


## MDS in Sklearn

MDS is implemented in Sklearn


$$
\begin{aligned}
& \text { 8isisis }
\end{aligned}
$$

Measuring distance between two persistence diagrams


data1



Distance between PD(data1) and PD(data2)

## Bottleneck distance between two persistent diagrams

- Given two persistence diagrams $X$ and $Y$, let $\eta$ be a bijection between points in the diagram. The following two distances are commonly used in the context of PH to measure the distance between two persistence diagrams:

$$
\begin{gathered}
W_{\infty}(X, Y)=\inf _{\eta: X \rightarrow Y} \sup _{x \in X}\|x-\eta(x)\|_{\infty} \\
W_{q}(X, Y)=\left[\inf _{\eta: X \rightarrow Y} \Sigma_{x \in X}\|x-\eta(x)\|_{\infty}^{q}\right]^{1 / q}
\end{gathered}
$$

Bottleneck distance between two persistent diagrams
MDSCALE with M


Matrix M describes the pair-wise distance
between the persistence $\rightarrow$ diagrams of each data element

MDS plot of the matrix $M$ with labels corresponding to each class.

##    100日时 $\rightarrow$ 8isesis sis 



Input data $\qquad$

## Remarks

- The axes obtained when drawing the MDS coordinates are , in themselves, meaningless.


## Remarks

- The axes obtained when drawing the MDS coordinates are, in themselves, meaningless.
- The orientation of the result MDS "picture" is arbitrary.


## Remarks

- The axes obtained when drawing the MDS coordinates are , in themselves, meaningless.
- The orientation of the result MDS "picture" is arbitrary.
- When we obtain MDS coordinates that have non-zero stress, we should remember that the distances among the resulting items are distorted representations of the relationships given by the input data. This distortion is greater when the stress is greater.
- That being said, we, in general, can rely on the larger distances as being more accurate than smaller distances.


## Non-Matric MDS

- Sometimes, there is no defined metric on points and all we are given is a similarity measure between the points.


## Non-Matric MDS

- Sometimes, there is no defined metric on points and all we are given is a similarity measure between the points.

The main idea in non-metric MDS :

- The actual values given to us are not that meaningful
- Ranking among different points is important
- Non-metric MDS finds a low-dimensional representation, which respects the ranking of distances as much as possible


## Non-Matric MDS

- Recall that in MDS we seek to find an optimal configuration xi that gives $d_{i j} \approx d_{i j}^{\prime}=\| x_{i}-x_{j}| |$ as close as possible.


## Non-Matric MDS

- Recall that in MDS we seek to find an optimal configuration xi that gives $d_{i j} \approx d_{i j}^{\prime}=\| x_{i}-x_{j}| |$ as close as possible.
- Relaxing $d i j \approx d^{\prime} i j$ from MDS by allowing $d^{\prime} i j \approx f\left(d_{i j}\right)$, for some monotone function $f$

Monotonic means : $d_{i j}<d_{k l} \Leftrightarrow f(d i j) \leq f\left(d_{k l}\right)$

## Non-Matric MDS

- Recall that in MDS we seek to find an optimal configuration xi that gives $d_{i j} \approx d_{i j}^{\prime}=\| x_{i}-x_{j}| |$ as close as possible.
- Relaxing $d i j \approx d^{\prime} i j$ from MDS by allowing $d^{\prime} i j \approx f\left(d_{i j}\right)$, for some monotone function $f$

Monotonic means : $d_{i j}<d_{k l} \Leftrightarrow f(d i j) \leq f\left(d_{k l}\right)$

Given a dimension d, non-metric MDS seeks to find an optimal configuration $X \subset R^{d}$ that gives $f\left(d_{i j}\right) \approx d^{\wedge} i j=\left\|x_{i}-x_{j}\right\|$ as close as possible.

- $f\left(d_{i j}\right)=d^{*} i j$ is only required to preserve the order of $d_{i j}$,
i.e., $d_{i j}<d_{k l} \Leftrightarrow f\left(d_{i j}\right) \leq f\left(d_{k l}\right) \Leftrightarrow d_{i j}^{*} \leq d_{k l}^{*}$


## Non-Matric MDS

The stress function for non-metric MDS is given by :

$$
\text { Stress }=\left(\sum_{i<j}\left(\hat{d}_{i j}-f\left(d_{i j}\right)\right)^{2} / \sum d_{i j}^{2}\right)^{\frac{1}{2}}
$$

Non-metric MDS optimizes over both position of the points of points and $f$

## Non-Matric MDS

The stress function for non-metric MDS is given by :

$$
\text { Stress }=\left(\sum_{i<j}\left(\hat{d}_{i j}-f\left(d_{i j}\right)\right)^{2} / \sum d_{i j}^{2}\right)^{\frac{1}{2}}
$$

Non-metric MDS optimizes over both position of the points of points and $f$

Solved numerically using (isotonic regression); we usually use classical MDS as starting initial position.

## Non-Matric MDS

The stress function for non-metric MDS is given by :

$$
\text { Stress }=\left(\sum_{i<j}\left(\hat{d}_{i j}-f\left(d_{i j}\right)\right)^{2} / \sum d_{i j}^{2}\right)^{\frac{1}{2}}
$$

Non-metric MDS optimizes over both position of the points of points and $f$

Solved numerically using (isotonic regression); we usually use classical MDS as starting initial position.


## ISOMAP

- Isomap extends MDS by utilizing geodesic distances induced by some neighborhood graph.
- There is no essential difference between MDS and ISOMAP algorithm once we find the distance matrix induced by the neighborhood graph.
- Isomap extends MDS by utilizing geodesic distances induced by some neighborhood graph.
- There is no essential difference between MDS and ISOMAP algorithm once we find the distance matrix induced by the neighborhood graph.

So the steps for ISOMAP on a given data :
1- Construct the neighborhood graph of the data $X$ using one of the neighborhood graphs we studied earlier in the course

- Isomap extends MDS by utilizing geodesic distances induced by some neighborhood graph.
- There is no essential difference between MDS and ISOMAP algorithm once we find the distance matrix induced by the neighborhood graph.

So the steps for ISOMAP on a given data :
1- Construct the neighborhood graph of the data $X$ using one of the neighborhood graphs we studied earlier in the course
2- Use the Dijekstra algorithm or the Floyd-Warshall algorithm to find the distance between nodes on the graph

- Isomap extends MDS by utilizing geodesic distances induced by some neighborhood graph.
- There is no essential difference between MDS and ISOMAP algorithm once we find the distance matrix induced by the neighborhood graph.

So the steps for ISOMAP on a given data :
1- Construct the neighborhood graph of the data $X$ using one of the neighborhood graphs we studied earlier in the course
2- Use the Dijekstra algorithm or the Floyd-Warshall algorithm to find the distance between nodes on the graph
3- Apply MDS on the distance matrix above and extract the coordinates with the desired dimension

## ISOMAP

A


B


C


A- Euclidian distance might not represent the actual distance between the points in the data.
B- We can construct the neighborhood graph of the data and then compute the geodesic distance between the points of the graph
C- Embedding the space we obtained in B into the plane.

## Appendix : Floyd-Warshall algorithm

```
let dist be a |V| x |V| array of minimum distances
initialized to infinity
for each edge (u,v)
    dist[u][v] \leftarroww(u,v) // the weight of the edge (u,v)
for each vertex v
    dist[V][V]}\leftarrow
for k from 1 to |V|
    for i from 1 to |V|
        for j from 1 to |V|
            if dist[i][j] > dist[i][k] + dist[k][j]
                        dist[i][j] \leftarrow dist[i][k] + dist[k][j]
                end if
```

Floyd algorithm is good to use when we want to compute the distance matrix on a dense graph. When the graph G is sparse, Dijekstra algorithm is a better choice.

